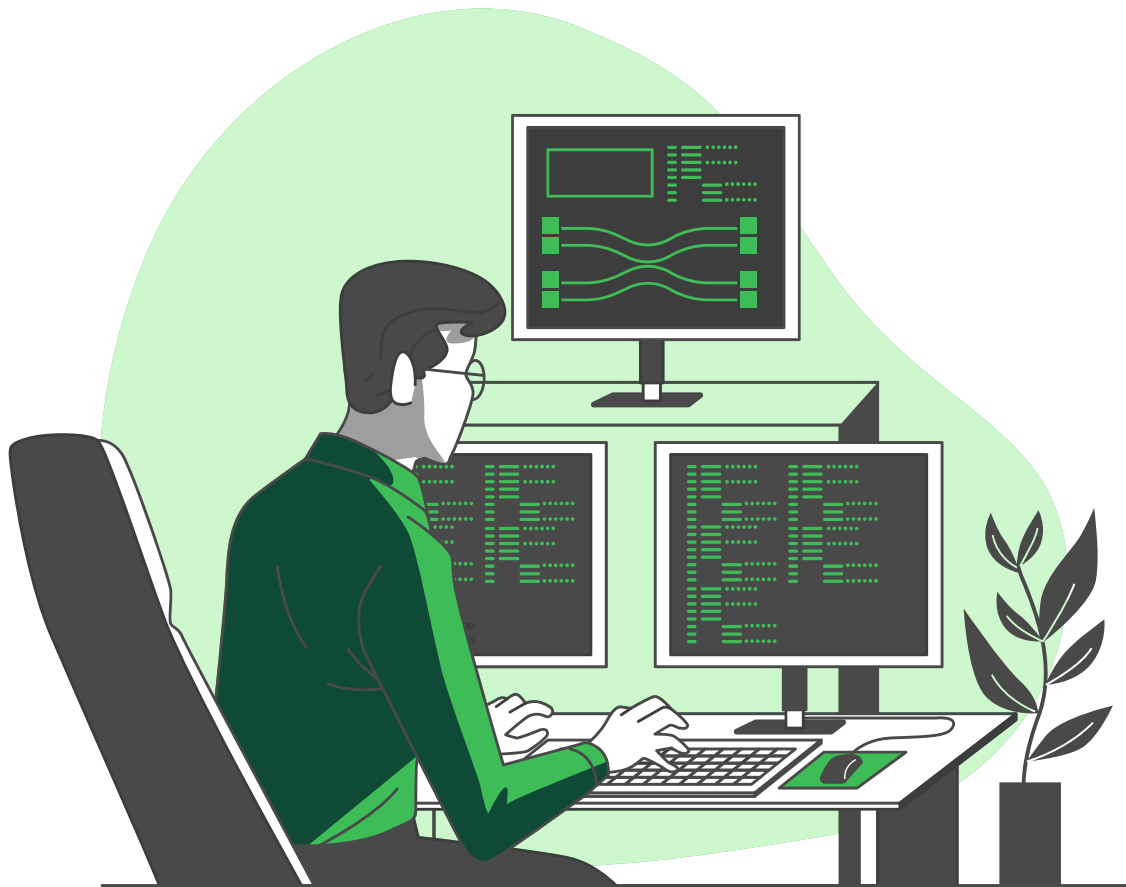
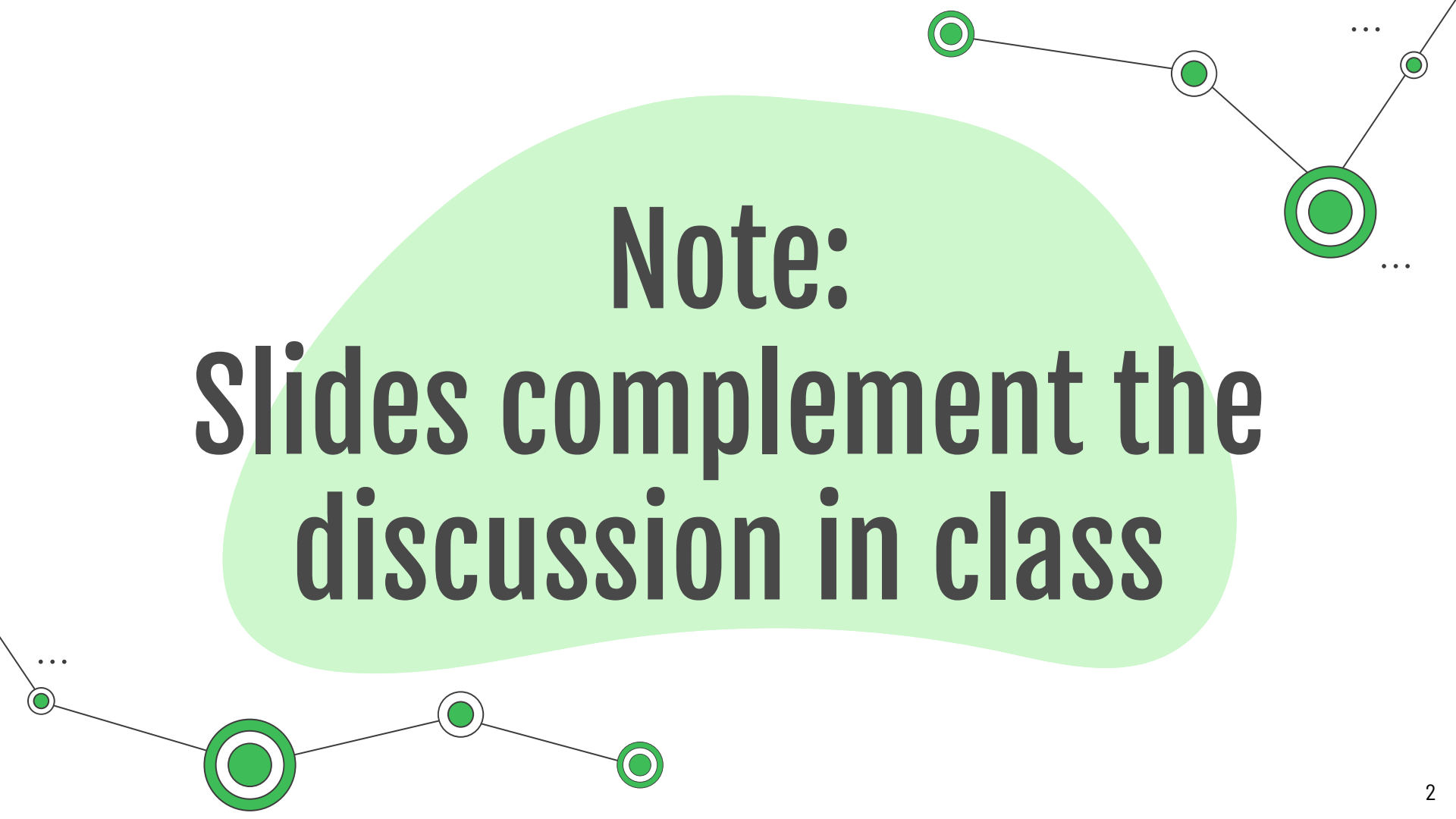




Graph Traversals

CS 251 - Data Structures
and Algorithms

A decorative network diagram consisting of several green circular nodes connected by thin black lines. Some nodes are single green circles, while others are double green circles. The nodes are arranged in a non-linear fashion, with some at the top right and others at the bottom left. Ellipses (...) are placed near some nodes to indicate a larger, unseen network.

Note:
**Slides complement the
discussion in class**

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Subgraphs & Trees

They are also graphs

02

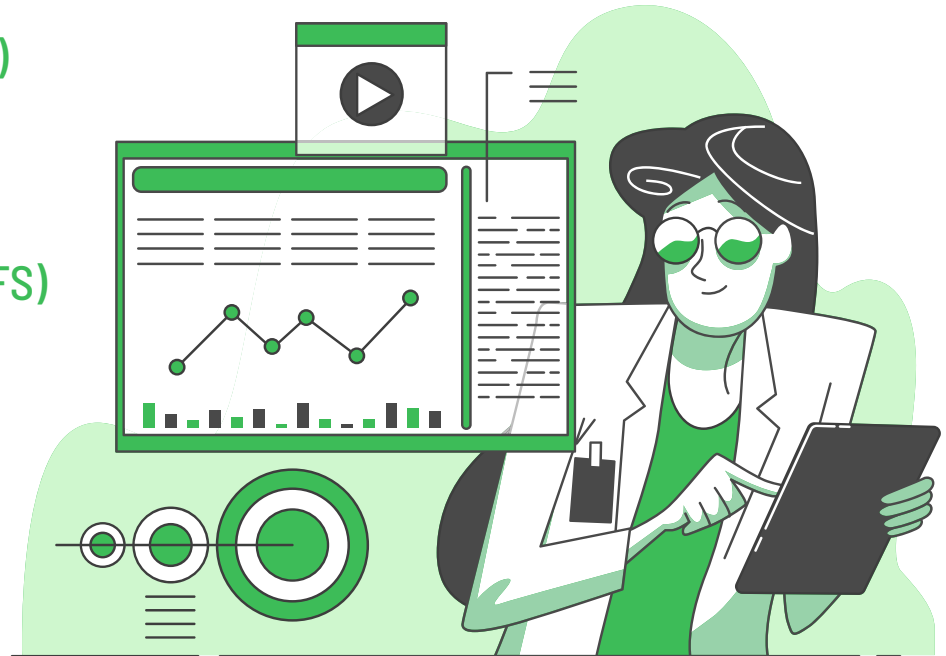
Depth-First Search (DFS)

Traverse by branches

03

Breadth-First Search (BFS)

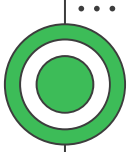
Traverse by levels



01

Subgraphs & Trees

They are also graphs



...



...



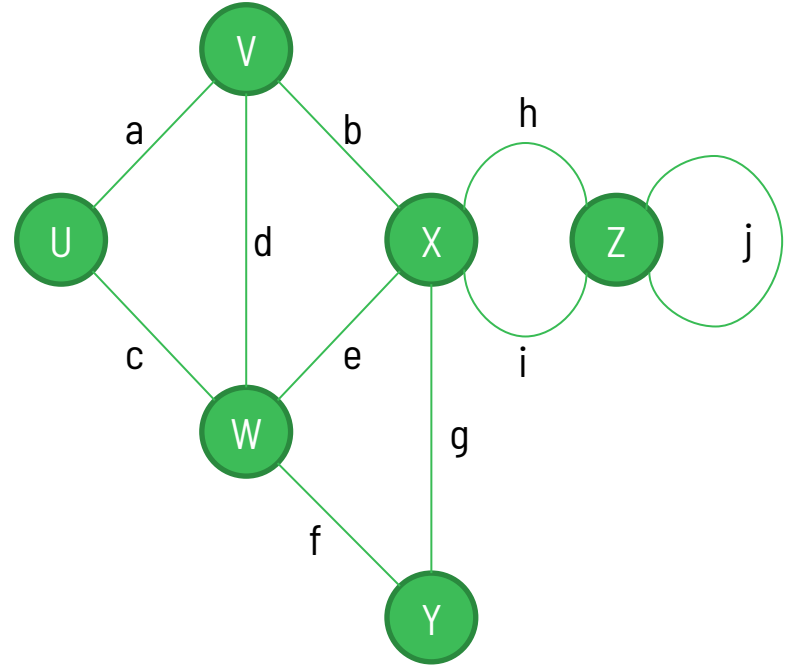
...

Revisit Subgraphs

Subgraph (of a graph): a graph made up of a subset of the vertices of G and a subset of the edges of G .

E.g.,

$$S = \{\{U, V, W, X, Y\}, \{a, b, c, d, e, f, g\}\}$$

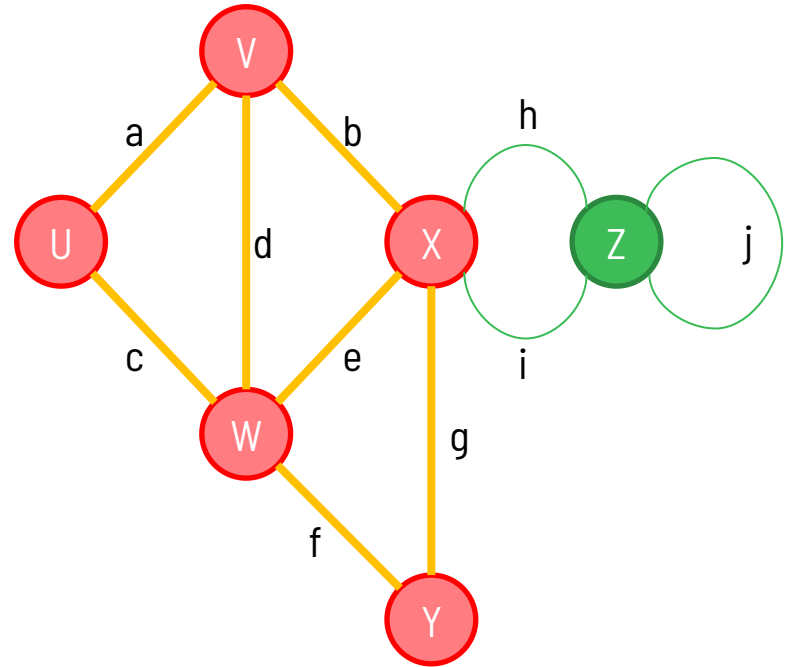


Revisit Subgraphs

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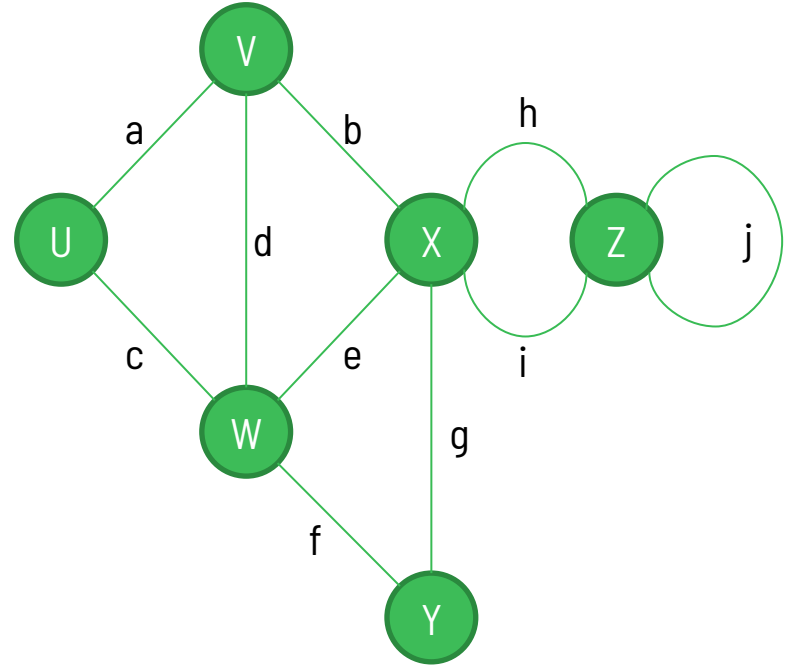


Revisit Subgraphs

Spanning Subgraph (of a graph): a subgraph that includes all vertices of G .

E.g.,

$S = \{\{U, V, W, X, Y, Z\}, \{a, b, c, d, e, g, h\}\}$

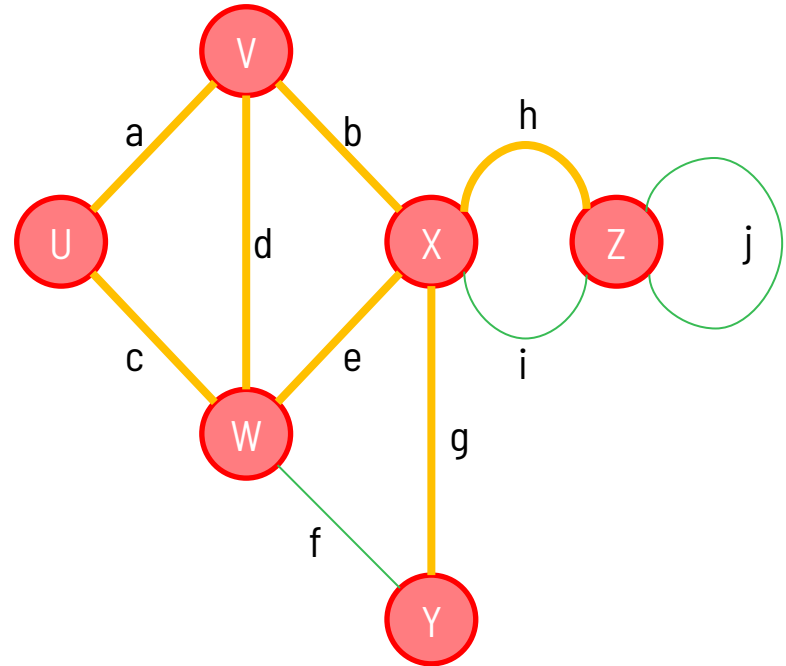


Revisit Subgraphs

Spanning Subgraph (of a graph): a subgraph that includes all vertices of G .

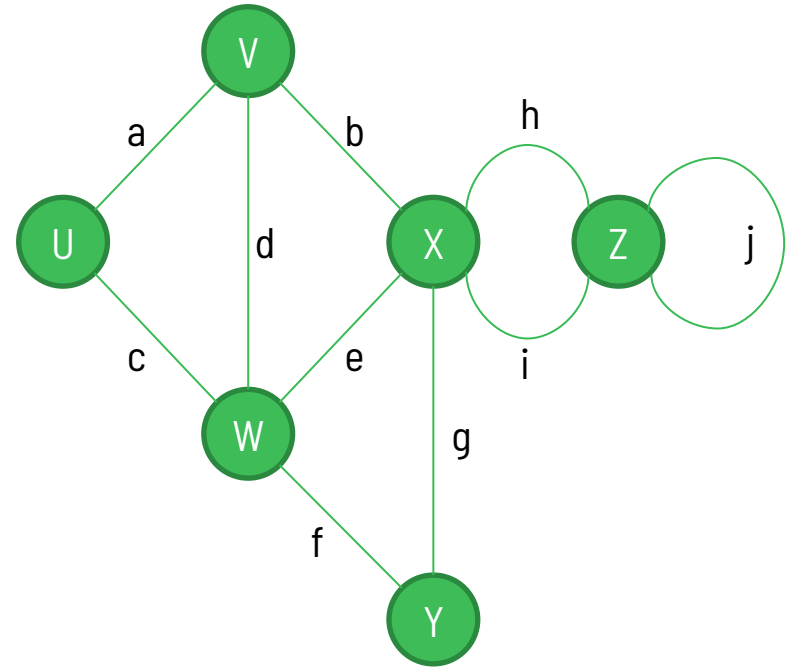
E.g.,

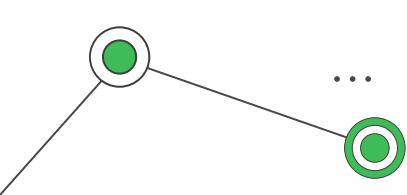
$S = \{\{U, V, W, X, Y, Z\}, \{a, b, c, d, e, g, h\}\}$



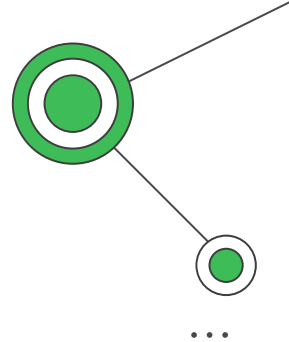
Connected Graphs

A **connected graph** is a graph where there is a path from every vertex to every other vertex.



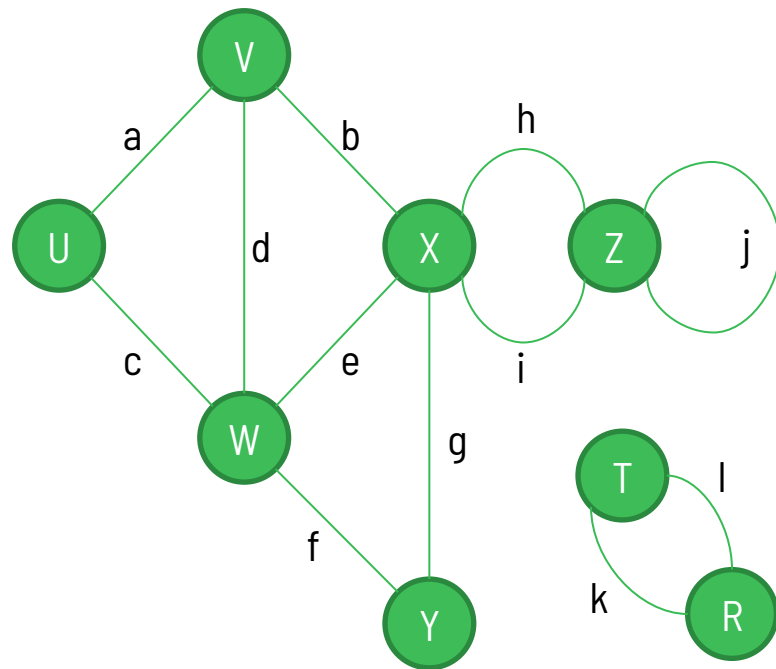


Connected Graphs



A **connected graph** is a graph where there is a path from every vertex to every other vertex.

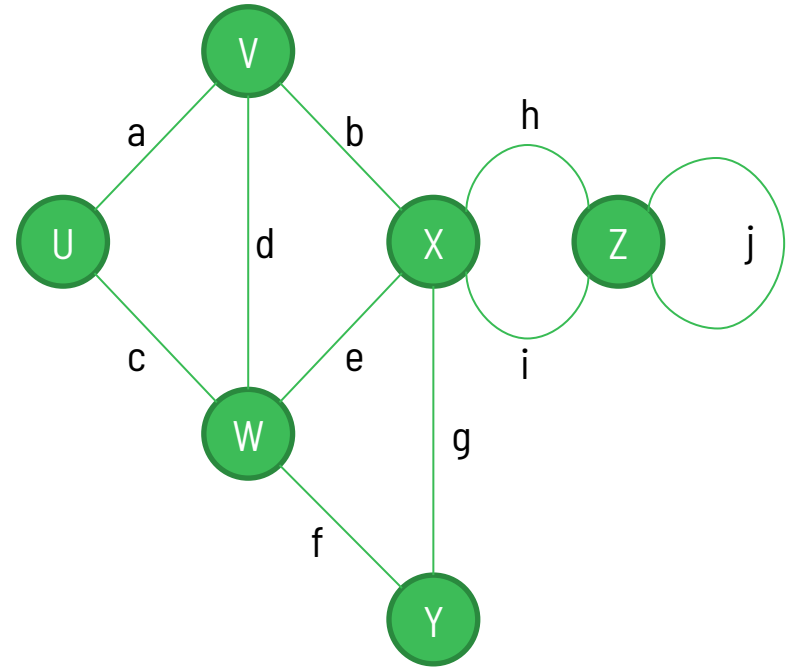
A **connected component** of a graph G is a maximal connected subgraph of G .



Let's Revisit Trees

A **tree** is an undirected graph T such that:

- T is connected
- T has no cycles (acyclic)



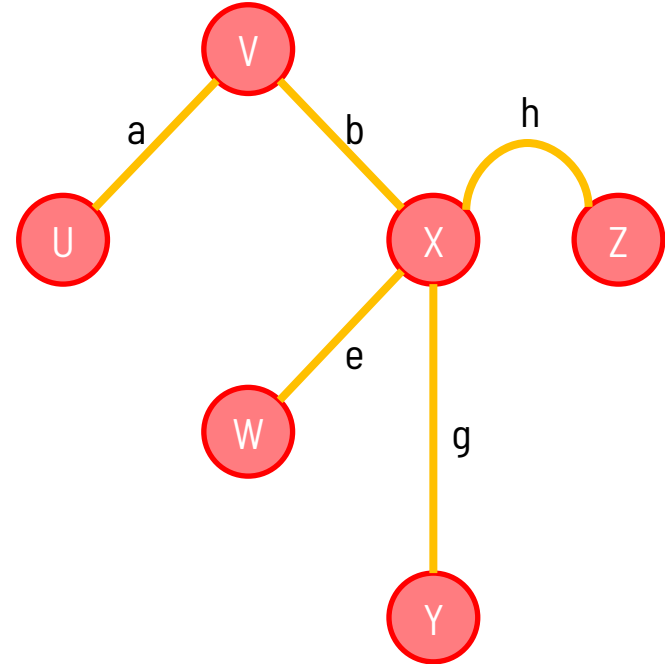
Let's Revisit Trees

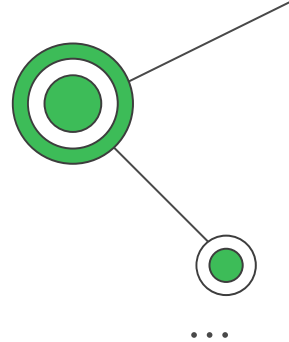
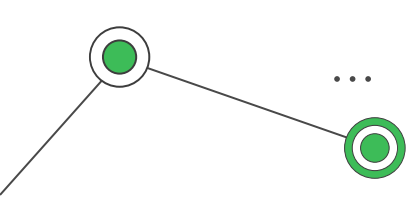
A **tree** is an undirected graph T such that:

- T is connected
- T has no cycles (acyclic)

A forest is an undirected graph without cycles (i.e., one or more trees).

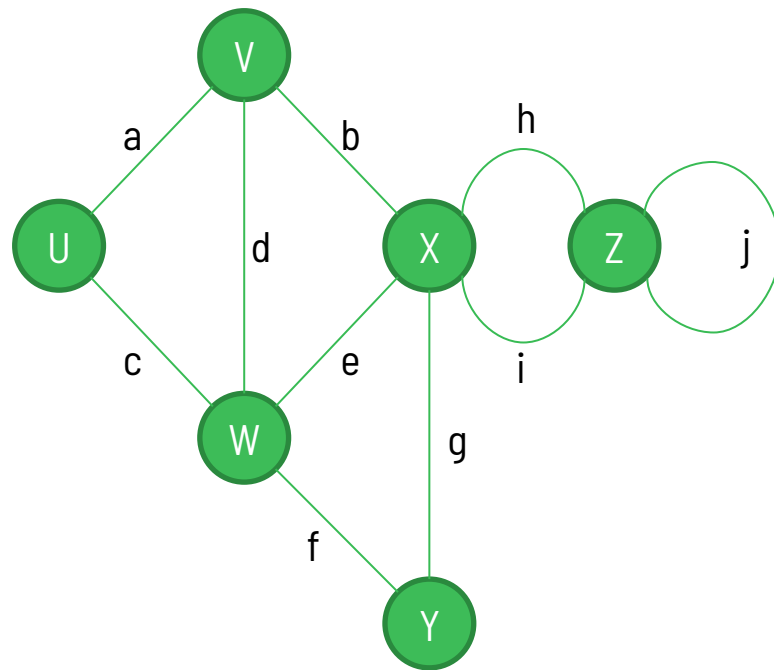
The connected components of a forest are trees.

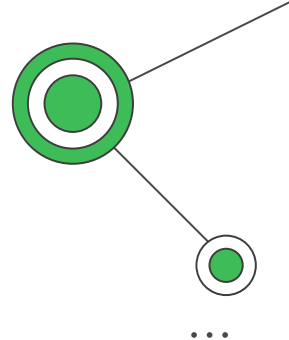
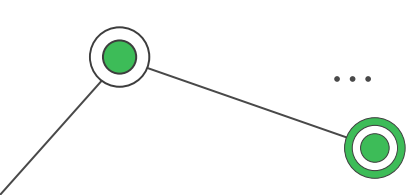




Spanning Tree

A **spanning tree** of a connected graph G is a subgraph of G that is a tree.





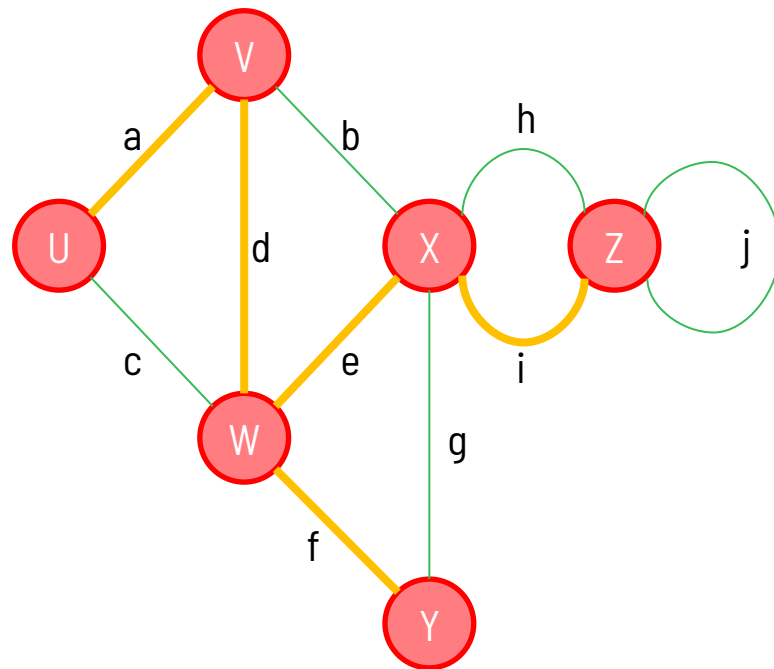
Spanning Tree

A **spanning tree** of a connected graph G is a subgraph of G that is a tree.

Q: Given a graph G , is it possible to have more than one spanning tree?

A: Yes, unless G is already a tree.

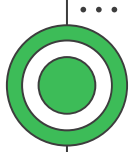
A **spanning forest** is a subgraph that is a forest.



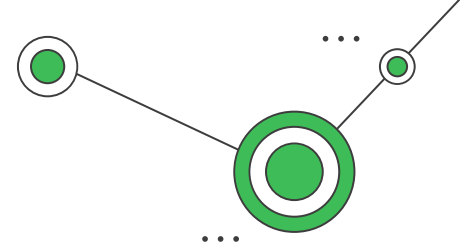
02

Depth-First Search (DFS)

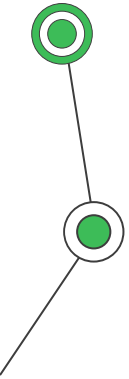
Traverse by branches

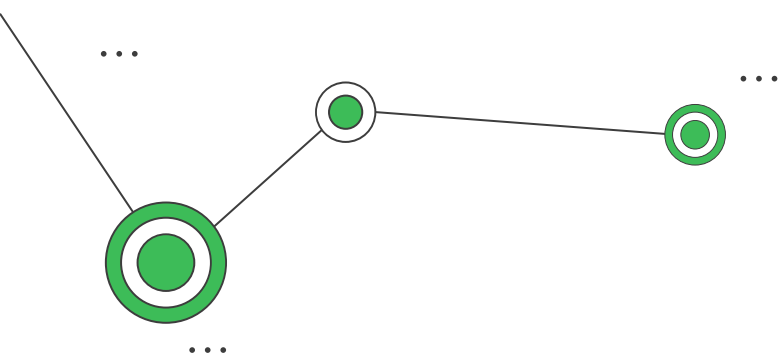


Depth-First Search (DFS)

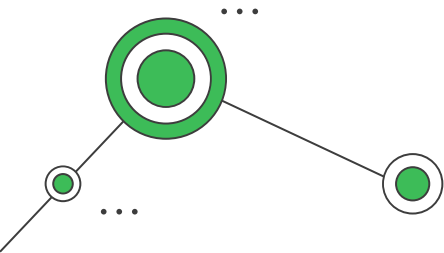


- A general technique for traversing a graph.
- Visits all the vertices and edges of a graph.
- Determines whether a graph is connected or not.
- Computes the connected components of a graph.
- Computes a spanning forest of a graph (spanning tree if the graph is connected).
- Useful as basis for solving other problems (e.g., finding a path between two vertices).





Recursive DFS



```
algorithm DFS( $G(V,E)$ ,  $s \in V$ ,  $T$ )
```

```
    mark  $s$  as visited
```

```
    for each  $w \in V$  adjacent to  $s$  do
```

```
        if  $w$  is not visited then
```

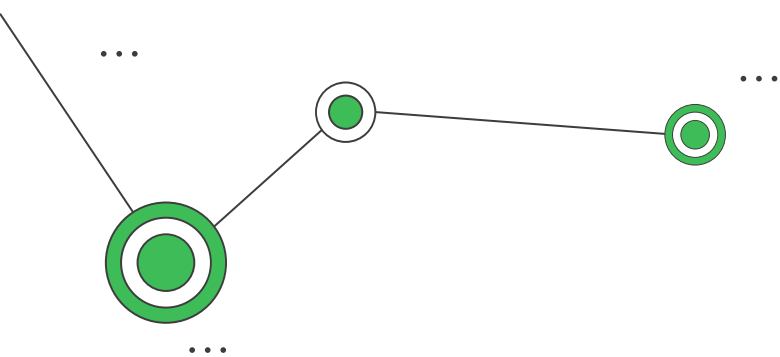
```
            insert  $(s,w)$  to  $T$ 
```

```
            DFS( $G$ ,  $w$ )
```

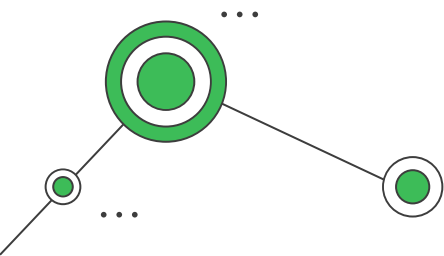
```
        end if
```

```
    end for
```

```
end algorithm
```



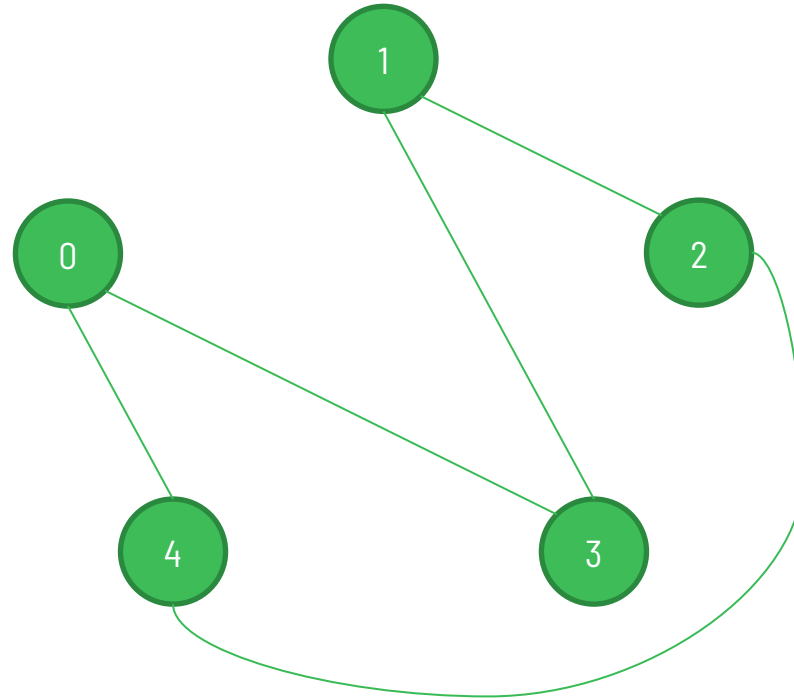
Iterative DFS



```
algorithm DFS( $G(V,E)$ ,  $s \in V$ )  
  let  $S$  be an empty stack  
   $S.push(s)$   
  
  while  $S$  is not empty do  
     $u \leftarrow S.pop()$   
    if  $u$  is not visited then  
      mark  $u$  as visited  
      for each vertex  $w$  adjacent to  $u$  do  
         $S.push(w)$   
      end for  
    end if  
  end while  
  
end algorithm
```

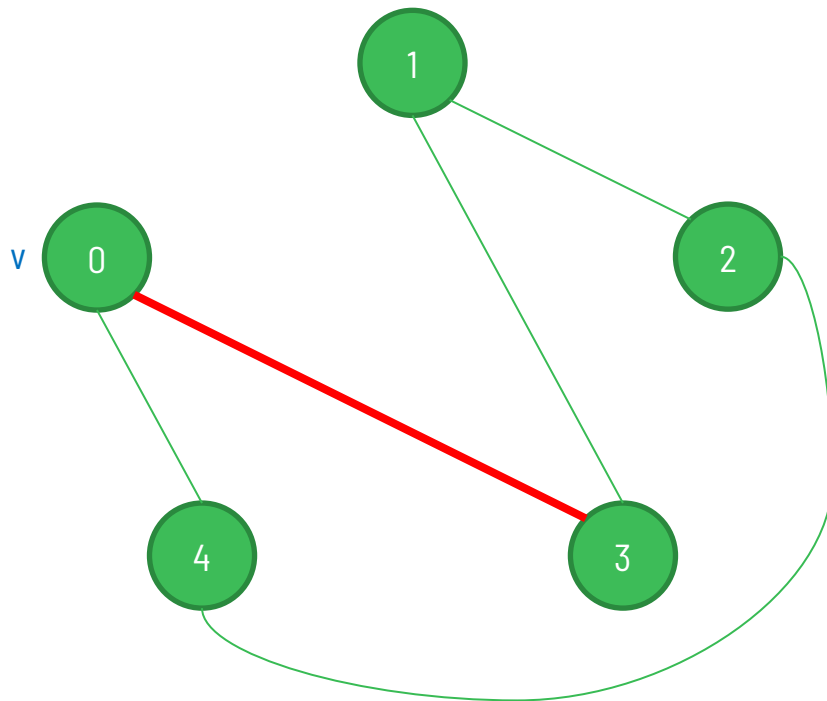
Show the DFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}



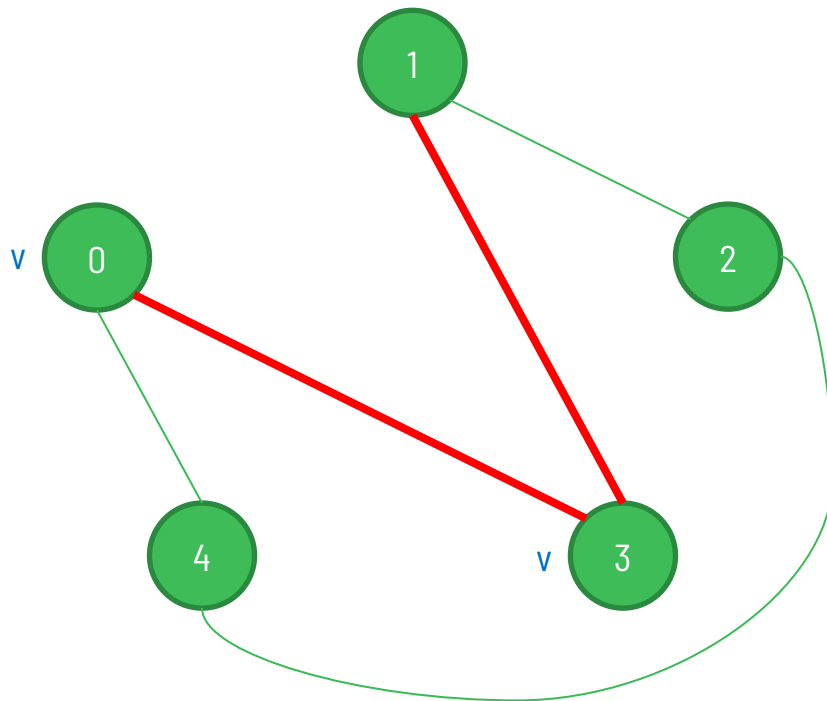
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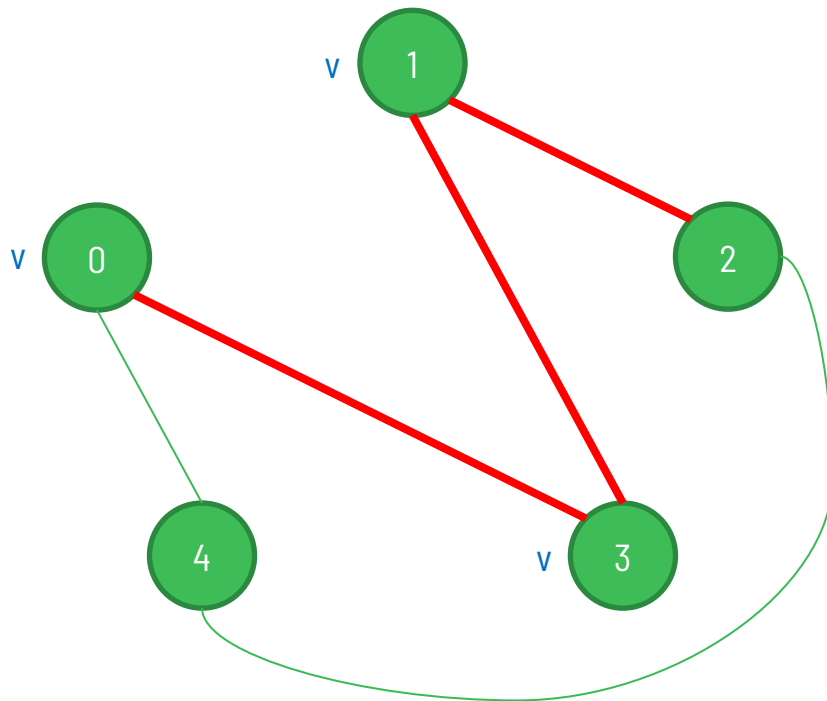
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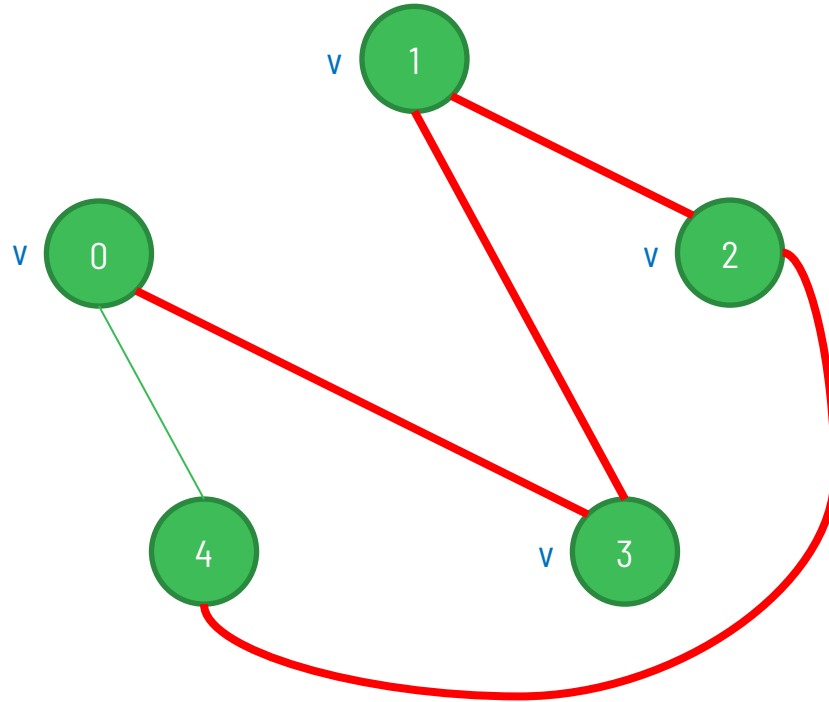
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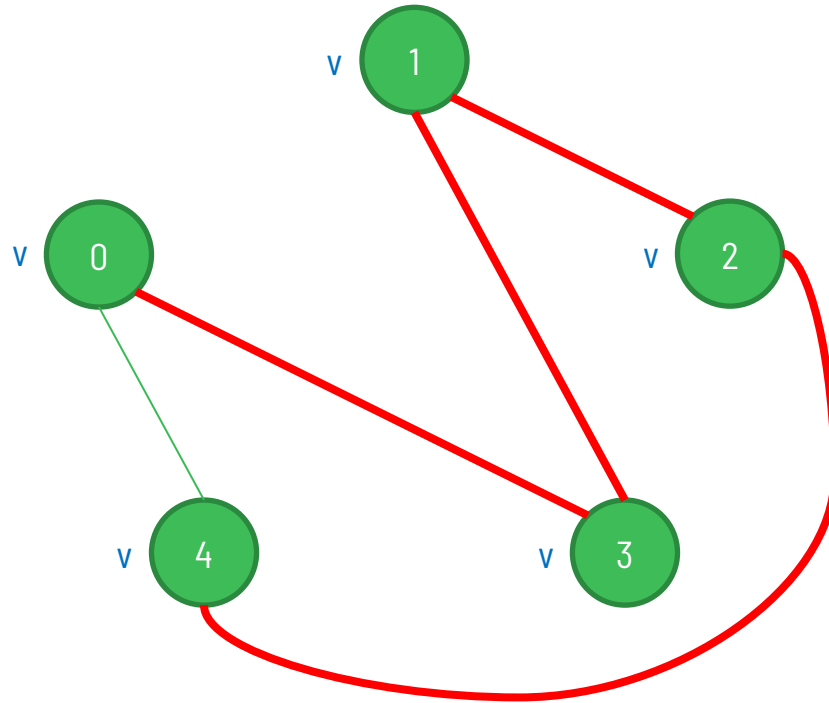
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2:	{1, 4}
3:	{0, 1}
4:	{0, 2}

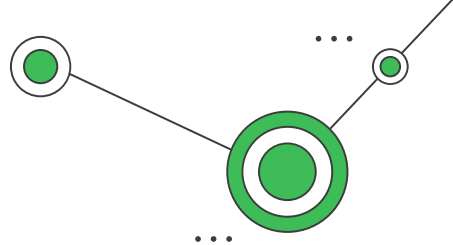


Show the DFS algorithm on the following undirected graph, starting at vertex 0.

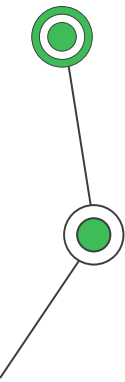
0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}



DFS Analysis



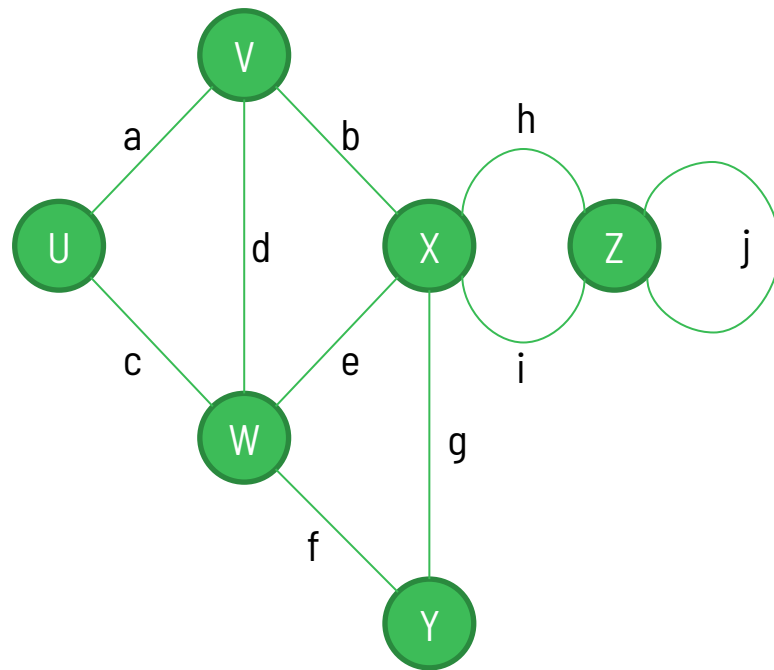
- By the end of the algorithm, we visited all the nodes ($|V|$).
- We traverse every edge twice ($2|E|$).
- Total runtime with adjacency list representation: $O(|V| + |E|)$.
- What happens if DFS ends, and we got at least one non-visited node?



Application: Finding a Path

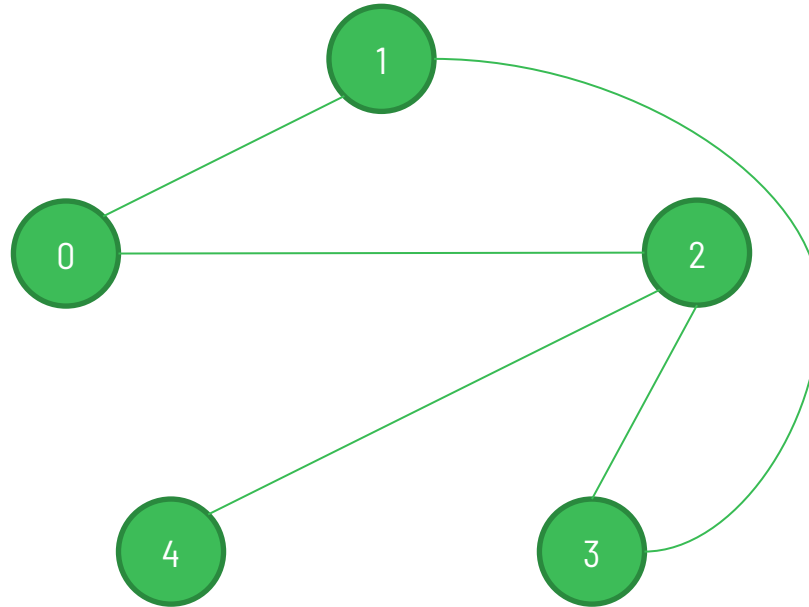
Given a graph $G(V, E)$, a source vertex $v \in V$, and a destination vertex $w \in V$, determine if there is a path from v to w , and if there is, find the path.

Solution: Use $\text{DFS}(G, v)$ and a stack S to keep track of the current path.



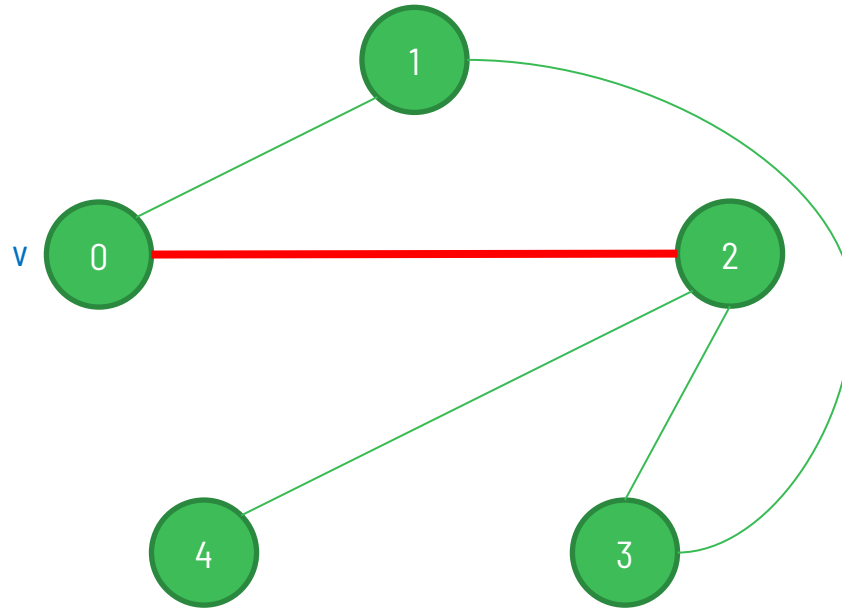
Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

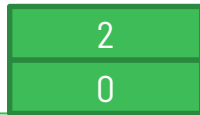
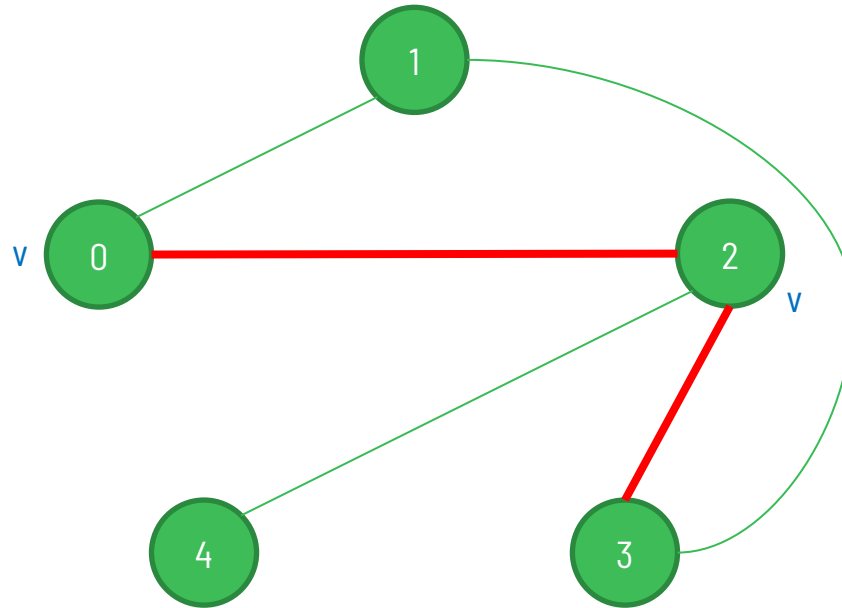
0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



0

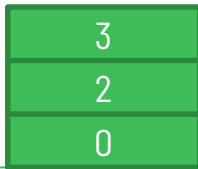
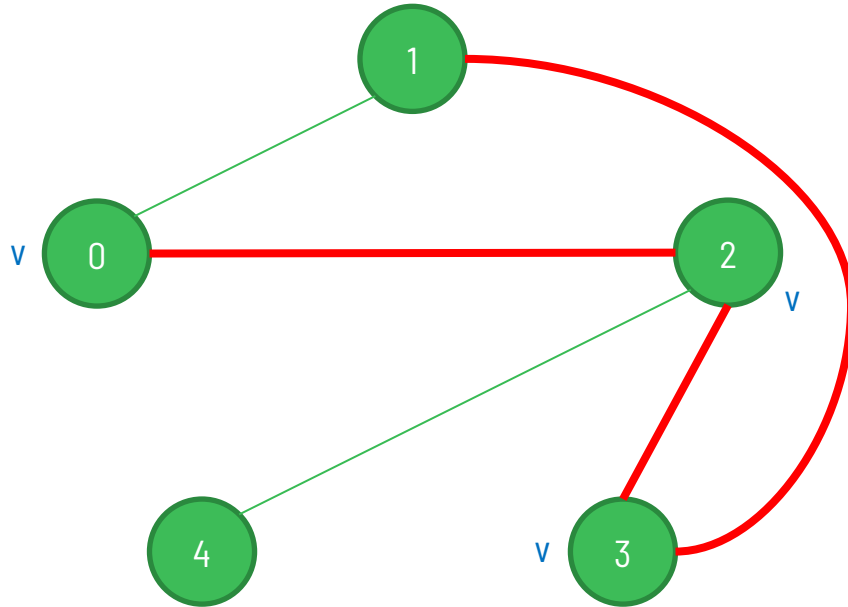
Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



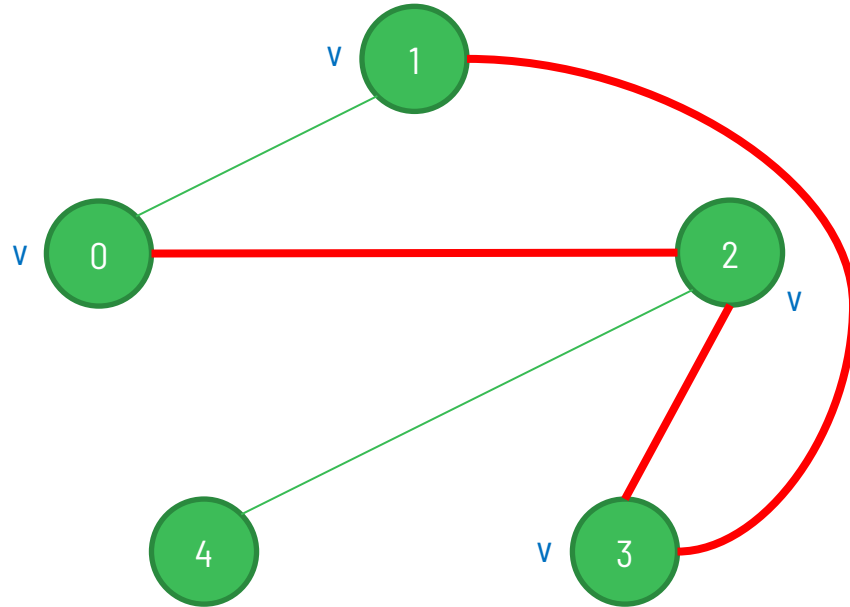
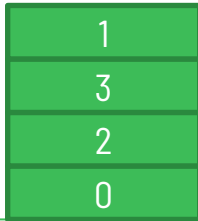
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0:	{2, 1}
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2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



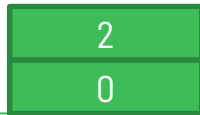
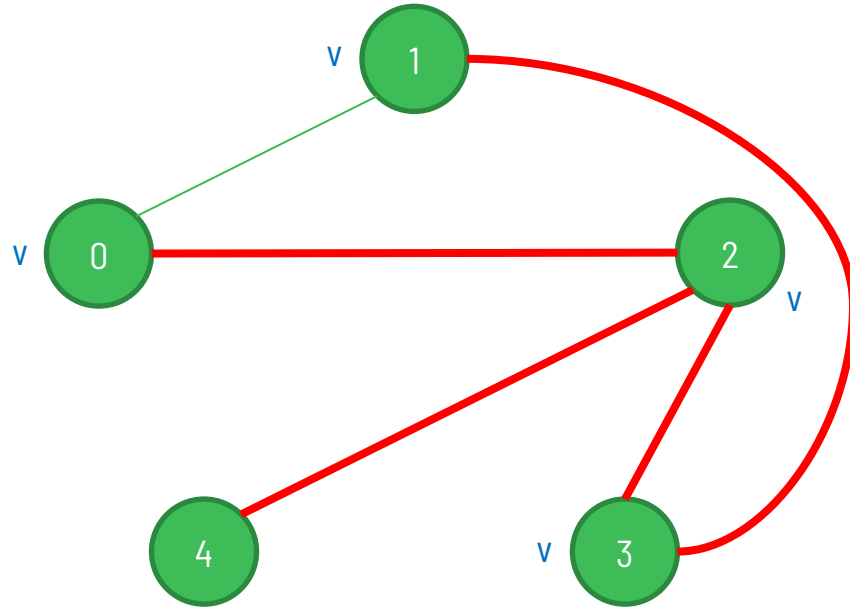
Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



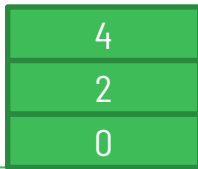
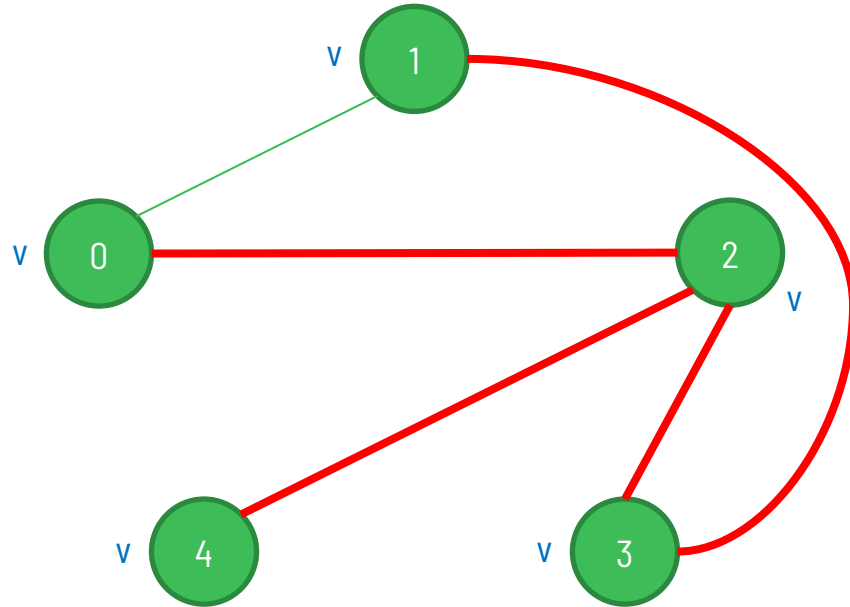
Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



Use DFS to find a path from vertex 0 to vertex 4 in the undirected graph below:

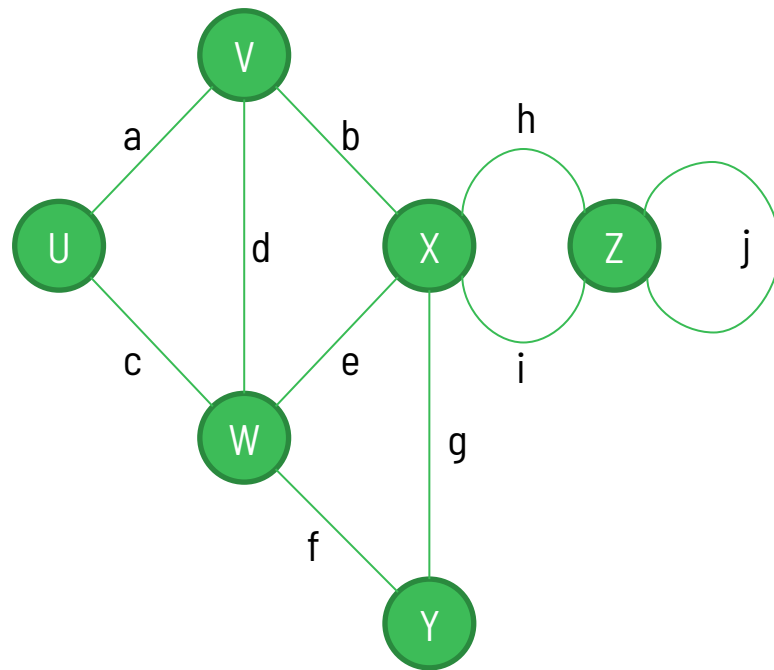
0:	{2, 1}
1:	{0, 3}
2:	{0, 3, 4}
3:	{1, 2}
4:	{2}



Application: Finding a Simple Cycle

Given a graph $G(V, E)$, find a simple cycle.

Solution: Use $\text{DFS}(G, v)$ and a stack S to keep track of the current path. Repeat until the DFS algorithm is complete, or until you encounter a back edge. Then, return the stack from that edge back to the first occurrence of the repeated vertex.



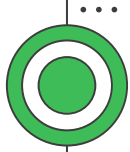
Fun Application: Mazes



03

Breadth-First Search (BFS)

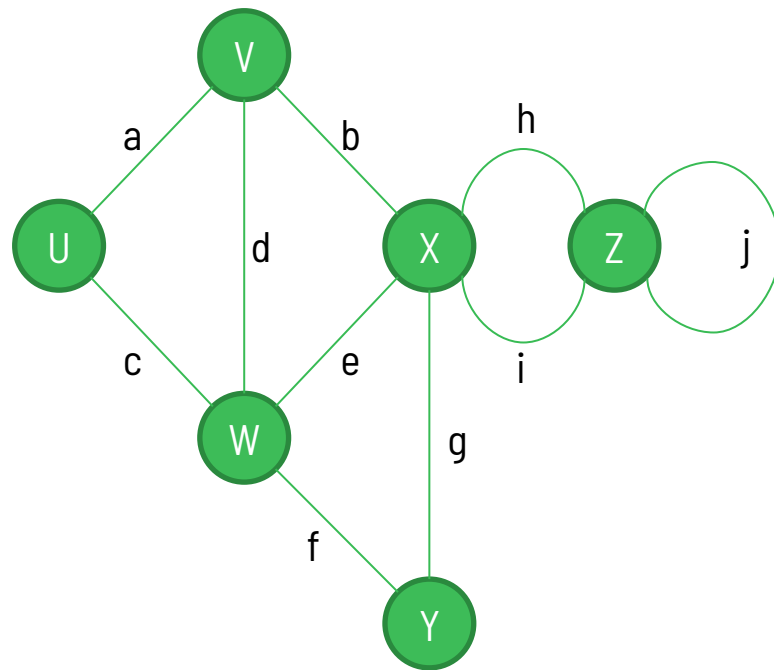
Traverse by levels



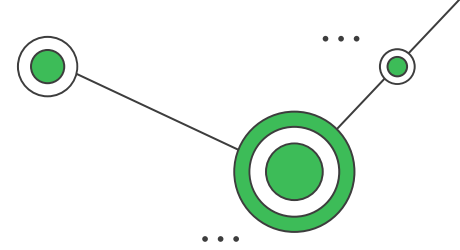
Single-Source Shortest Path

Given a graph $G(V, E)$ and a source vertex $v \in V$, find the shortest path (if a path exists) to a target vertex $w \in V$.

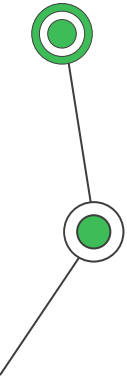
Can we do this with DFS?

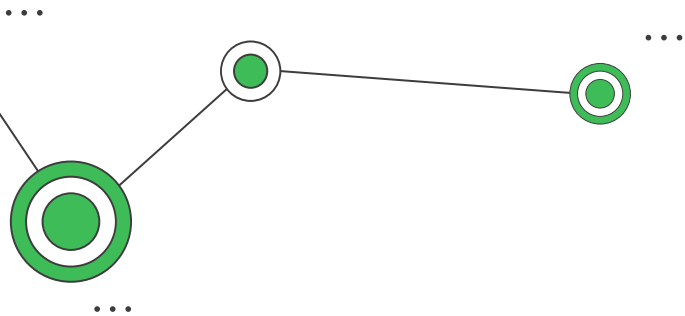


Breadth-First Search (BFS)

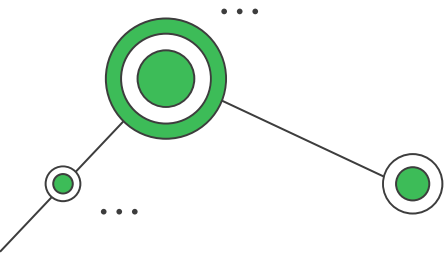


- A general technique for traversing a graph.
- Visits all the vertices and edges of a graph.
- It computes the distance (smallest number of edges) from a single vertex to each reachable vertex.
- For any vertex v reachable from u , the simple path in the BFS tree from u to v corresponds to a "shortest path" from u to v in the graph.





Breadth-First Search (BFS)



```
algorithm BFS( $G(V,E)$ ,  $s \in V$ )
```

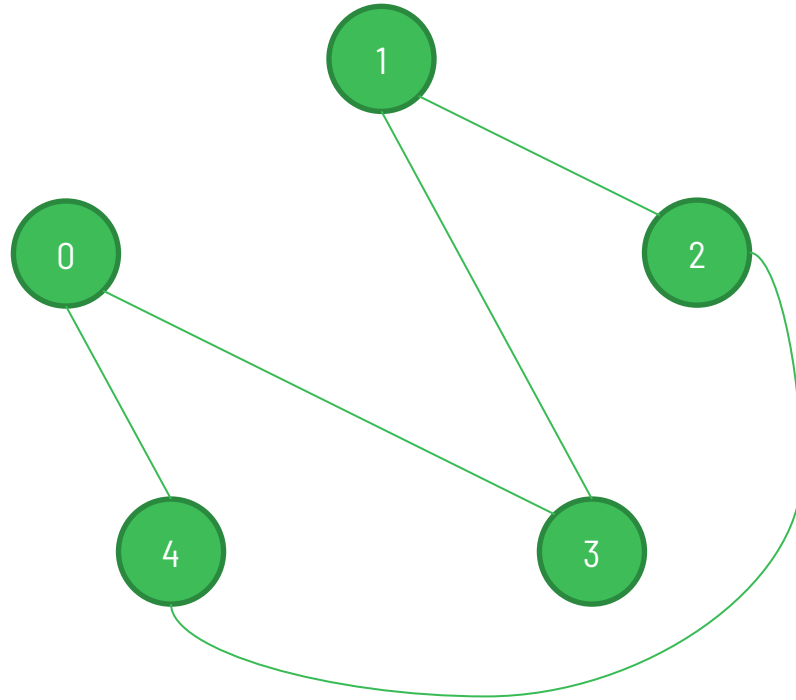
```
    let  $Q$  be an empty queue  
    let edgeTo be an array of size  $|V|$   
    mark  $s$  as visited  
     $Q.enqueue(s)$ 
```

```
    while  $Q$  is not empty do  
         $u \leftarrow Q.dequeue()$   
        for all  $w$  adjacent to  $u$  do  
            if  $w$  is not visited then  
                mark  $w$  as visited  
                 $Q.enqueue(w)$   
                edgeTo[ $w$ ]  $\leftarrow u$   
            end if  
        end for  
    end while
```

```
    return edgeTo  
end algorithm
```

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}



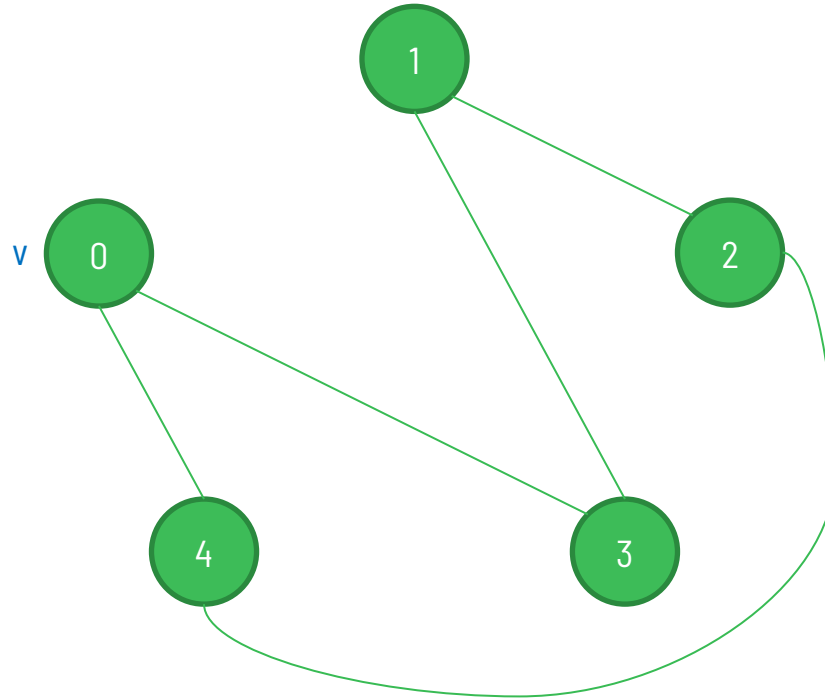
Q: {}

edgeTo

0	1	2	3	4

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}



Q: {0}

edgeTo

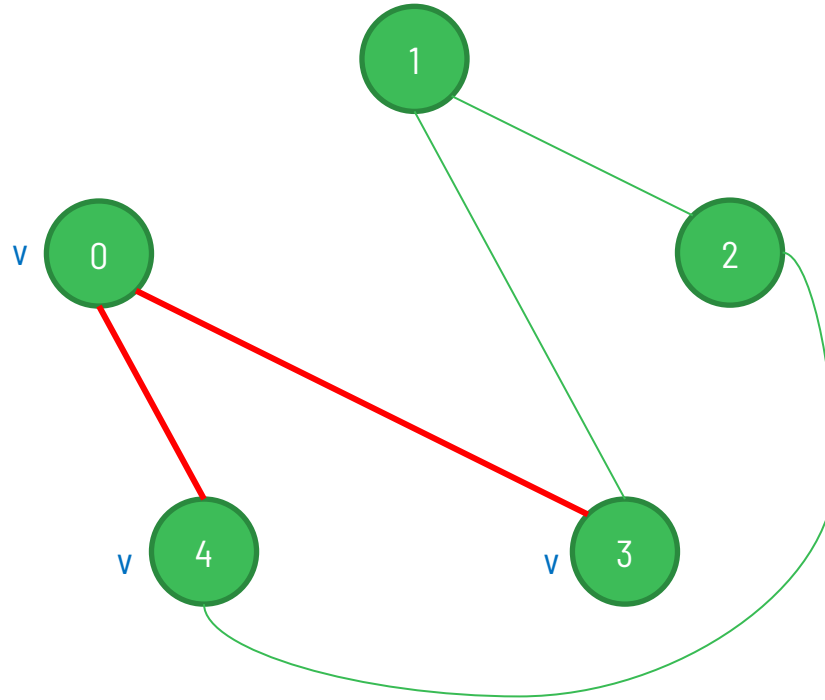
0	1	2	3	4

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}

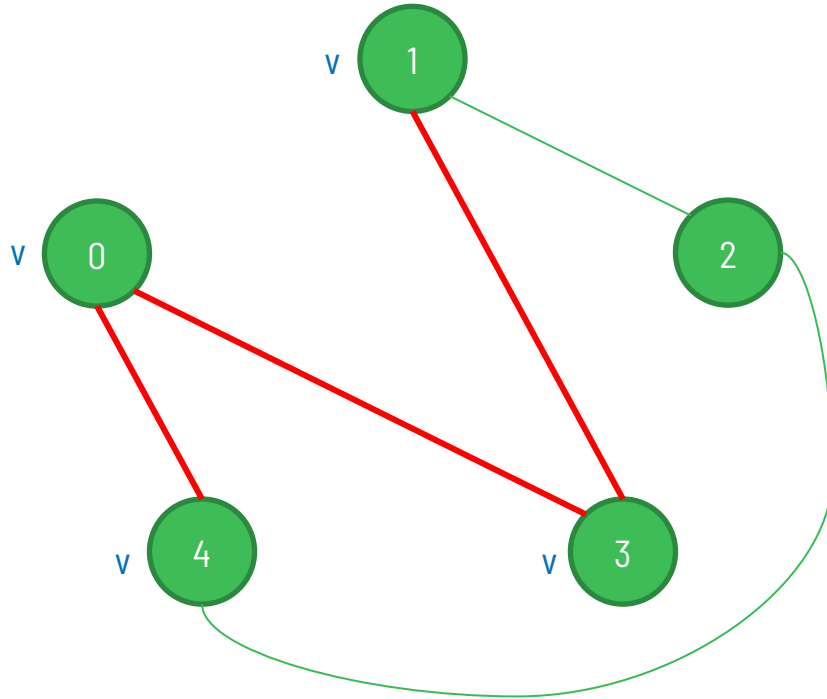
Q: {3, 4}

	0	1	2	3	4
edgeTo				0	0



Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}

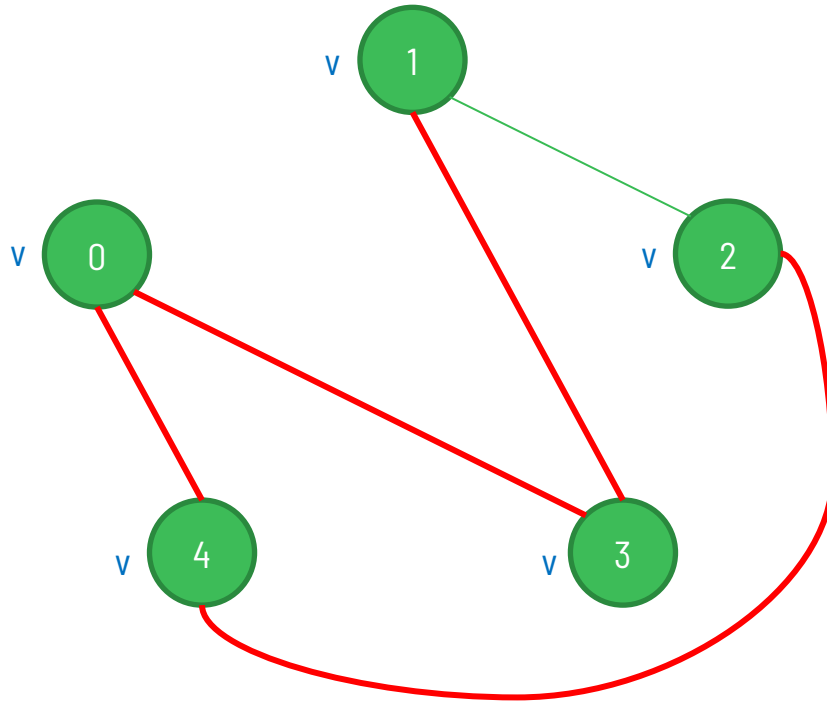


Q: {4, 1}

	0	1	2	3	4
edgeTo		3		0	0

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}

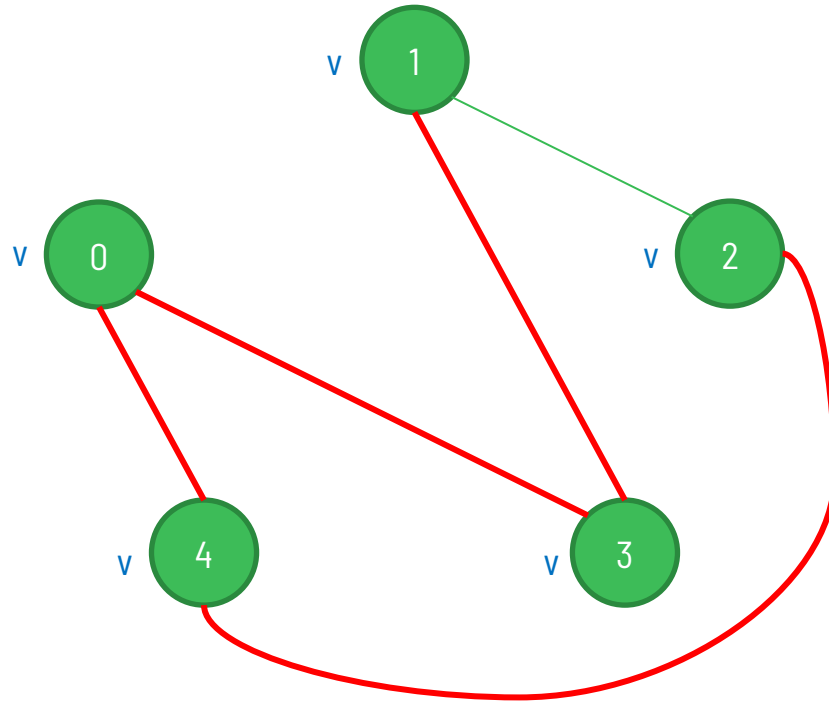


Q: {1, 2}

	0	1	2	3	4
edgeTo		3	4	0	0

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}



Q: {2}

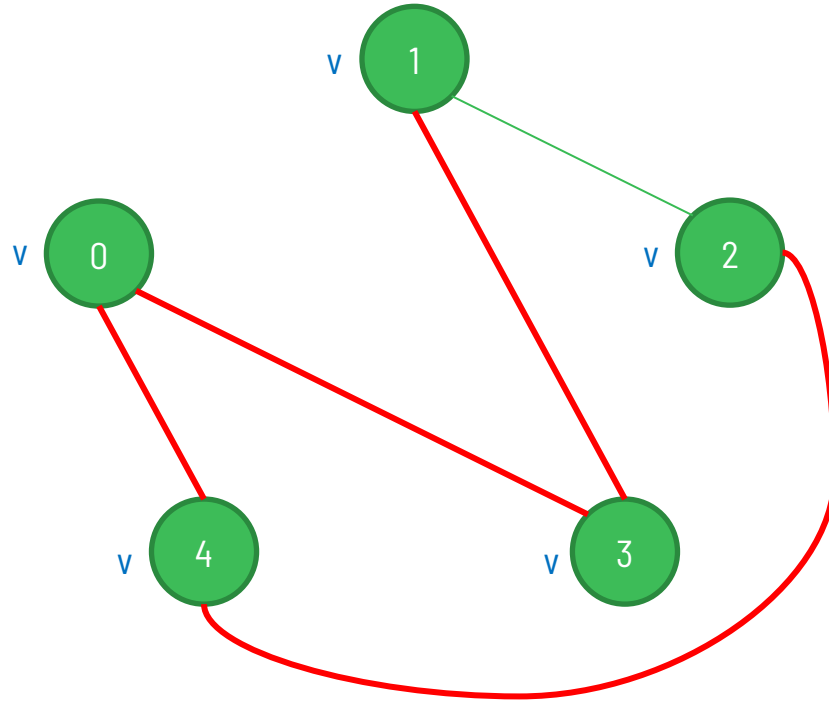
	0	1	2	3	4
edgeTo		3	4	0	0

Show the BFS algorithm on the following undirected graph, starting at vertex 0.

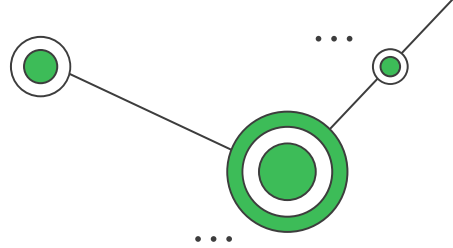
0:	{3, 4}
1:	{2, 3}
2:	{1, 4}
3:	{0, 1}
4:	{0, 2}

Q: {}

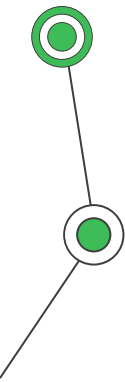
	0	1	2	3	4
edgeTo		3	4	0	0



BFS Analysis



- By the end of the algorithm, we visited all the nodes ($|V|$).
- We traversed every edge twice ($2|E|$)
- Total runtime with adjacency list representation: $O(|V| + |E|)$
- edgeTo array contains the traversal of the graph.



DFS -vs- BFS

DFS

LIFO style

Goes to the end of a path before coming back to an intersection

BFS

FIFO style

Tries multiple paths, one edge at a time

Next Slide = $\emptyset$

Do you have any questions?

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